

A New Algorithm to Obtain the Adjugate Matrix using CUBLAS on GPU

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Abstract—In this paper a parallel code for obtain the Adjugate Matrix with real coefficients are used. We postulate a new linear transformation in matrix product form and we apply this linear transformation to an augmented matrix $(A|I)$ by means of both a minimum and a complete pivoting strategies, then we obtain the Adjugate matrix. Furthermore, if we apply this new linear transformation and the above pivot strategy to a augmented matrix $(A|b)$, we obtain a Cramer's solution of the linear system of equations. That new algorithm present an $O(n^3)$ computational complexity when $(A,b) \in \mathbb{R}^n$. We use subroutines of CUBLAS 2^{nd} and 3^{rd} levels in double precision and we obtain correct numeric results.

Index Terms—Adjoint matrix, Adjugate matrix, Cramer's rule, CUBLAS, GPU.

I. INTRODUCTION

A Linear System of Equations (LSE) can be defined as a set of m equations with n unknowns represented by a matrix A , a vector b and an unknown vector x , namely, $Ax = b$. Many methods have been proposed to solve such linear equations. A famous one is Cramer's rule, where each component of the solution is determined as the ratio of two determinants.

When trying to solve a system of n equations using Cramer's rule, one needs to compute $n+1$ determinants, each of order n . If these are computed in a straightforward way, using the Laplace Expansion, the solution to the linear system takes $(n+1)n!(n-1)$ multiplications, plus a similar number of additions. Although Cramer's rule possesses a fundamental theoretical importance, it may result impractical in computations. It is for that reason that this method is seldom recommended [1]-[6]. Cramer's rule has at least one attractive property: it computes every element of the solutions independently. For this reason, it can be a practical method for some special linear systems on parallel computers [7].

Another approach, with a certain mathematical appeal but considerable computational pitfalls, finds the solution to a linear system of equations using the inverse matrix A^{-1} . However, in virtually every application, it is unnecessary and inadvisable to compute the inverse matrix explicit. The inverse requires more arithmetic and produces a less accurate answer. Therefore, neither of the above methods are recommended [8]. Again, in this paper, we will propose a new efficient method to calculate the adjugate matrix.

The Gaussian Transformation (GT) for solving an LSE has

proved to be the best option for most practical applications. The new transformation proposed here can be obtained from it. Next, we briefly review this topic.

II. LU-MATRICAL DECOMPOSITION WITH GT. NOTATION AND DEFINITIONS

The problem of solving a linear system of equations $Ax = b$ is central to the field of matrix computation. There are several ways to perform the elimination process necessary for its matrix triangulation. We will focus on the Doolittle-Gauss elimination method: the algorithm of choice when A is square, dense, and un-structured.

Let us assume that $A \in \mathbb{R}^{mn}$ is nonsingular and that we wish to solve the linear system $Ax = b$. Here we show how for exact arithmetic and partial pivoting and column interchanges some Gauss transformations M_1, \dots, M_{n-1} can almost always be found such that $M_{n-1}, \dots, M_2 M_1 A = U$ is upper triangular [9]. The original $Ax = b$ problem is then equivalent to the upper triangular system $Ux = (M_{n-1}, \dots, M_2 M_1) b$ which can be solved through back-substitution.

Suppose, then, that $A \in \mathbb{R}^{mn}$ and that, for some $k < n$, we have determined the Gauss transformations

$M_1, \dots, M_{k-1} \in \mathbb{R}^{mn}$ such that

$$A^{(k-1)} \equiv M_{k-1} \cdots M_1 A = \begin{pmatrix} A_{11}^{(k-1)} & A_{12}^{(k-1)} \\ \mathbf{0} & A_{22}^{(k-1)} \end{pmatrix} \begin{matrix} (k-1) \\ (n-k+1) \end{matrix}$$

Where: $A_{11}^{(k-1)}$ is an upper triangular matrix?

$$\text{Now, if } A_{22}^{(k-1)} = \begin{pmatrix} a_{kk}^{(k-1)} & \cdot & \cdot & a_{kn}^{(k-1)} \\ \cdot & & & \cdot \\ \cdot & & & \cdot \\ \cdot & & & \cdot \\ a_{nk}^{(k-1)} & \cdot & \cdot & a_{mm}^{(k-1)} \end{pmatrix}$$

and $a_{kk}^{(k-1)} \neq 0$, then the *multiplicators* :

$$m_i = \frac{a_{ik}^{(k-1)}}{a_{kk}^{(k-1)}}; i = k+1, \dots, n; \text{ with } a_{kk} \neq 0 \text{ are well defined.}$$

So, we have the following

$$U = MA = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-2} \\ 0 & \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{vmatrix} & \dots & \dots & a_{n,n-2} \\ 0 & 0 & \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} & \dots & \dots & a_{n-2,n-2} \end{pmatrix};$$

$$m_{n-1,2} = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-2} \\ a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-2} \\ a_{3,1} & a_{3,2} & a_{3,3} & \dots & \dots & a_{3,n-2} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-2,1} & a_{n-2,2} & a_{n-2,3} & \dots & \dots & a_{n-2,n-2} \end{pmatrix};$$

n
 $k=1,2,\dots,n-1$

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-1} & a_{1,n} \\ a_{21} & a_{22} & a_{23} & \dots & \dots & a_{2,n-1} & a_{2,n} \\ a_{31} & a_{32} & a_{33} & \dots & \dots & a_{3,n-1} & a_{3,n} \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \dots & \dots & a_{n-1,n-1} & a_{n-1,n} \\ a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-1} & a_{n,n} \end{pmatrix};$$

$$m_{n-1,3} = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-2} \\ a_{21} & a_{22} & a_{23} & \dots & \dots & a_{2,n-2} \\ a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-2} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-2,1} & a_{n-2,2} & a_{n-2,3} & \dots & \dots & a_{n-2,n-2} \end{pmatrix};$$

$$M = M'_{n-1} \dots M'_1 =$$

$$= \begin{pmatrix} m_{11} & 0 & 0 & \dots & \dots & 0 & 0 \\ m_{21} & m_{22} & 0 & \dots & \dots & 0 & 0 \\ m_{31} & m_{32} & m_{33} & \dots & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ m_{n-1,1} & m_{n-1,2} & m_{n-1,3} & \dots & \dots & m_{n-1,n-1} & 0 \\ m_{n,1} & m_{n,2} & m_{n,3} & \dots & \dots & m_{n,n-1} & m_{n,n} \end{pmatrix}$$

$$m_{n-1,n-1} = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-2} \\ a_{21} & a_{22} & a_{23} & \dots & \dots & a_{2,n-2} \\ a_{31} & a_{32} & a_{33} & \dots & \dots & a_{3,n-2} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-2,1} & a_{n-2,2} & a_{n-2,3} & \dots & \dots & a_{n-2,n-2} \end{pmatrix}$$

where

$$m_{11} = 1 ; m_{21} = -a_{21} ; m_{22} = a_{11} ;$$

$$m_{31} = -\begin{vmatrix} a_{31} & a_{32} \\ a_{21} & a_{22} \end{vmatrix} ; m_{32} = -\begin{vmatrix} a_{11} & a_{12} \\ a_{31} & a_{32} \end{vmatrix} ; m_{33} = \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} ;$$

$$m_{n,1} = \begin{pmatrix} a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-1} \\ a_{21} & a_{22} & a_{23} & \dots & \dots & a_{2,n-1} \\ a_{31} & a_{32} & a_{33} & \dots & \dots & a_{3,n-1} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \dots & \dots & a_{n-1,n-1} \end{pmatrix};$$

$$m_{n-1,1} = \begin{pmatrix} a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-2} \\ a_{21} & a_{22} & a_{23} & \dots & \dots & a_{2,n-2} \\ a_{31} & a_{32} & a_{33} & \dots & \dots & a_{3,n-2} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-2,1} & a_{n-2,2} & a_{n-2,3} & \dots & \dots & a_{n-2,n-2} \end{pmatrix};$$

$$m_{n,2} = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \dots & \dots & a_{1,n-1} \\ a_{n,1} & a_{n,2} & a_{n,3} & \dots & \dots & a_{n,n-1} \\ a_{31} & a_{32} & a_{33} & \dots & \dots & a_{3,n-1} \\ \dots & \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \dots & \dots & a_{n-1,n-1} \end{pmatrix};$$

$$m_{n,3} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n-1} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n-1} \\ a_{n,1} & a_{n,2} & a_{n,3} & \cdot & \cdot & \cdot & a_{n,n-1} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \cdot & \cdot & \cdot & a_{n-1,n-1} \end{vmatrix};$$

$$u_{3,n} = \begin{vmatrix} a_{11} & a_{12} & a_{1,n} \\ a_{21} & a_{22} & a_{2,n} \\ a_{31} & a_{32} & a_{3,n} \end{vmatrix}$$

The Laplace Expansion of sub-matrix $A_{n-1,n-1}$ taking out the last column, is equivalent to multiply the $(n-1)$ -th row of the matrix M by the $(n-1)$ -th column of A . Then we have

$u_{n-1,n-1} = |A_{n-1,n-1}|$ and

$$m_{n,n-1} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n-1} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n-1} \\ a_{31} & a_{32} & a_{33} & \cdot & \cdot & \cdot & a_{3,n-1} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n,1} & a_{n,2} & a_{n,3} & \cdot & \cdot & \cdot & a_{n,n-1} \end{vmatrix};$$

$$u_{n-1,n-1} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n-1} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n-1} \\ a_{31} & a_{32} & a_{33} & \cdot & \cdot & \cdot & a_{3,n-1} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \cdot & \cdot & \cdot & a_{n-1,n-1} \end{vmatrix};$$

In a similar way, we have

$$m_{n,n} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n-1} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n-1} \\ a_{31} & a_{32} & a_{33} & \cdot & \cdot & \cdot & a_{3,n-1} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \cdot & \cdot & \cdot & a_{n-1,n-1} \end{vmatrix}$$

$$u_{n-1,n} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n} \\ a_{31} & a_{32} & a_{33} & \cdot & \cdot & \cdot & a_{3,n} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n-1,1} & a_{n-1,2} & a_{n-1,3} & \cdot & \cdot & \cdot & a_{n-1,n} \end{vmatrix};$$

and

$$U = MA = \begin{pmatrix} u_{11} & u_{12} & u_{13} & \cdot & \cdot & \cdot & u_{1,n-1} & u_{1,n} \\ 0 & u_{22} & u_{23} & \cdot & \cdot & \cdot & u_{2,n-1} & u_{2,n} \\ 0 & 0 & u_{33} & \cdot & \cdot & \cdot & u_{3,n-1} & u_{3,n} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & \cdot & \cdot & \cdot & u_{n-1,n-1} & u_{n-1,n} \\ 0 & 0 & 0 & \cdot & \cdot & \cdot & 0 & u_{n,n} \end{pmatrix}$$

Finally, taking the last row of M multiplied by the last column of A we have

$$u_{n,n} = \begin{vmatrix} a_{11} & a_{12} & a_{13} & \cdot & \cdot & \cdot & a_{1,n} \\ a_{21} & a_{22} & a_{23} & \cdot & \cdot & \cdot & a_{2,n} \\ a_{31} & a_{32} & a_{33} & \cdot & \cdot & \cdot & a_{3,n} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{n,1} & a_{n,2} & a_{n,3} & \cdot & \cdot & \cdot & a_{n,n} \end{vmatrix}$$

where

$$u_{11} = a_{11}; u_{12} = a_{12}; u_{13} = a_{13}; u_{1,n-1} = a_{1,n-1};$$

$$u_{1,n} = a_{1,n}$$

$$u_{22} = \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix}; u_{23} = \begin{vmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{vmatrix};$$

$$u_{2,n-1} = \begin{vmatrix} a_{11} & a_{1,n-1} \\ a_{21} & a_{2,n-1} \end{vmatrix}; u_{2,n} = \begin{vmatrix} a_{11} & a_{1,n} \\ a_{21} & a_{2,n} \end{vmatrix}$$

$$u_{33} = \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix}; u_{3,n-1} = \begin{vmatrix} a_{11} & a_{12} & a_{1,n-1} \\ a_{21} & a_{22} & a_{2,n-1} \\ a_{31} & a_{32} & a_{3,n-1} \end{vmatrix};$$

Now, if M is the new transformation, then we have

$$M = \prod_{k=1}^n M'_k; U = MA$$

In order to solve the linear system of equations $Ax = b$, we have

$$MAX = Mb$$

$$Ux = Mb$$

We can use the “backward process” and solve the linear system of equations using only determinants.

For a matrix A with floating point entries this process requires

$$\frac{n(n-1)(2n-1)}{6} + \frac{(n-2)(n-1)}{2} = \frac{n^3}{3} - \frac{4}{3}n + 1$$

Floating point multiplications [13].

Proof. Let it suffice to explain how the algorithm works with the following case.

Let $A \in \mathbb{I}^{n \times n}$ and

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}; P = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}; \Pi = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$$

$$M_{J_3} M_{J_2} M_{J_1} P A \Pi = |A| I$$

$$M_J = |A| (P A \Pi)^{-1}$$

$$M_J = \Pi^{-1} |A| A^{-1} P^{-1}$$

$$\Pi M_J P = A^{Adj}$$

For $k=1, \dots, 3$; $(a_{00}^{(1)} = 1)$ we have:

$$A^{Adj} = \Pi M_{J_3} M_{J_2} M_{J_1} P =$$

$$= \begin{pmatrix} \begin{vmatrix} a_{22} & a_{23} \\ a_{32} & a_{33} \end{vmatrix} & -\begin{vmatrix} a_{12} & a_{13} \\ a_{32} & a_{33} \end{vmatrix} & \begin{vmatrix} a_{12} & a_{13} \\ a_{22} & a_{23} \end{vmatrix} \\ -\begin{vmatrix} a_{21} & a_{23} \\ a_{31} & a_{33} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{13} \\ a_{31} & a_{33} \end{vmatrix} & -\begin{vmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{vmatrix} \\ \begin{vmatrix} a_{21} & a_{22} \\ a_{31} & a_{32} \end{vmatrix} & -\begin{vmatrix} a_{11} & a_{12} \\ a_{31} & a_{32} \end{vmatrix} & \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} \end{pmatrix}$$

A. Solution of $Ax=b$ with M_J

The simultaneous linear equations systems can also be solved with this new linear transformation.

$$M_J = \prod_{k=1}^n M_{J_k} = A^{Adj}$$

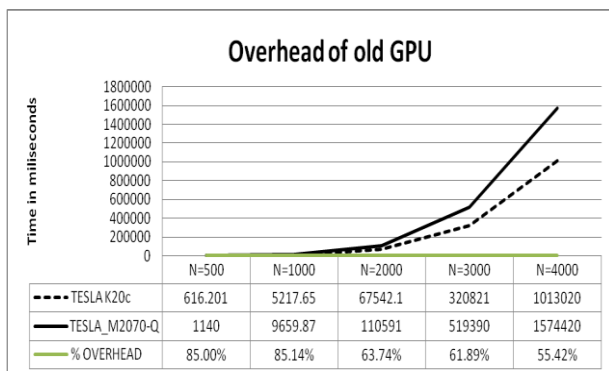
Since: , then:

$$A^{Adj} Ax = A^{Adj} b; ; |A|Ix = A^{Adj} b ; x = \frac{A^{Adj} b}{|A|I}$$

This result is a Cramer-type solution in $O(n^3)$. Consequently, this is a new result.

B. Numeric Results

We use two models of Tesla GPU: M2070-Q and K20c. We show the results in the following graph



V. CONCLUSION

In this paper we introduce a new theorem LU on the decomposition into determinants of matrix A and the new linear transformations, expressed as equations (1), (2), a

modified Doolittle-Gauss elimination process. Additionally we expressed the equations (3) and (4), a modified Doolittle-Gauss-Jordan elimination process to calculate the adjugate matrix.

Most simultaneous linear equation systems can also be solved with these new linear transformations. The result is Cramer-type solutions in $O(n^3)$. This fact is new.

We have also proposed a modified Doolittle-Gauss-Jordan elimination process in two versions: the first one applied to the A matrix and the second to the augmented matrix (A|b).

The first one is a new algorithm to compute determinants in exact form if and only if $A \in \mathbb{I}^n$, and the second is a new method to solve linear system of equations.

On the other hand, we have proposed a modified Doolittle-Gauss-Jordan elimination process in two version: the first one applied to the augmented matrix (A|I) and the second to the augmented matrix (A|b). The first version is a new algorithm to calculate, in exact form, the Adjugate matrix, if $A \in \mathbb{I}^n$. The second version is a new direct method to solve linear system of equations in exact form if $(A, b) \subset \mathbb{I}^n$. Provided that $(A, b) \subset \mathbb{R}^n$ the above algorithms calculate, in approximate form, Cramer-type solutions of the linear systems of equations.

Gaussian elimination is usually the most economical way to solve $Ax = b$. Nevertheless there are three reasons why this new method might be relevant when the matrix's coefficients are integers:

- (1) The flop counts tend to exaggerate the Gaussian elimination advantage.
- (2) The present method provides guaranteed stability; there is no "growth factor" to worry about as in Gaussian elimination.
- (3) In cases of ill-conditioning, the reliability of the present method is unsurpassed when $A \in \mathbb{I}^n$.

Finally, when referring to the Cramer's rule, it has been affirmed by G. Strang [16] that: "...Thus each component of x is a ratio of two determinants, a polynomial of degree n divided by another polynomial of degree n. This fact might have been recognized from Gauss elimination, but it never was". This is made evident in the present paper.

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