

# Alternative approach to simplex method for the solution of linear programming problem

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*Abstract- In this paper, new alternative methods for simplex method, Big M method and dual simplex method are introduced. These methods are easy to solve linear programming problem. These are powerful methods. It reduces number of iterations and save valuable time by skipping calculations of net evaluation.*

**Key words:** Linear programming problem, optimal solution, simplex method, alternative method.

## I. INTRODUCTION

Khobragade et al. [2, 3, 4] suggested an alternative approach to solve linear programming problem. In this paper, an attempt has been made to solve linear programming problem (LPP) by new method which is an alternative for simplex method. This method is different from Khobragade et al. [2-4] Method.

## II. AN ALTERNATIVE ALGORITHM FOR SIMPLEX METHOD

To find optimal solution of any LPP by an alternative method for simplex method, algorithm is given as follows:

Step 1. Check objective function of LPP is of maximization or minimization type. If it is to be minimization type then convert it into a maximization type by using the result:

$$\text{Min. } Z = - \text{Max. } (-Z).$$

Step 2. Check whether all  $b_i$  (RHS) are non-negative. If any  $b_i$  is negative then multiply the corresponding equation of the constraints by (-1).

Step 3. Express the given LPP in standard form then obtain initial basic feasible solution.

Step 4. Select  $\max \sum x_{ij}$ ,  $x_{ij} \geq 0$ , for entering vector.

Step 5. Choose greatest coefficient of decision variables.

(i) If greatest coefficient is unique, then element corresponding to this row and column becomes pivotal (leading) element.

(ii) If greatest coefficient is not unique, then use tie breaking technique.

Step 6. Use usual simplex method for this table and go to next step.

Step 7. Ignore corresponding row and column. Proceed to step 5 for remaining elements and repeat the same procedure until an optimal solution is obtained or there is an indication for unbounded solution.

Step 8. If all rows and columns are ignored, then current solution is an optimal solution.

### PROBLEM -1

$$\text{Max. } Z = 5x_1 + 3x_2$$

Subject to the constraints:

$$x_1 + x_2 \leq 2,$$

$$5x_1 + 2x_2 \leq 10,$$

$$3x_1 + 8x_2 \leq 12,$$

$$x_1, x_2 \geq 0.$$

SOLUTION: We have the constraints

$$x_1 + x_2 + s_1 = 2,$$

$$5x_1 + 2x_2 + s_2 = 10$$

$$3x_1 + 8x_2 + s_3 = 12$$

where  $S_1, S_2, S_3$  are slack variables.

### New Simplex Table

$C_B$	basis	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	2	1	1	1	0	0
0	$s_2$	10	5	2	0	1	0
0	$s_3$	12	3	8	0	0	1
0	$s_1$	1/2	5/8	0	1	0	-1/8
0	$s_2$	7	17/4	0	0	1	-1/4
3	$x_2$	3/2	3/8	1	0	0	1/8
0	$s_1$	-9/17	0	0	1	-5/34	-3/34
5	$x_1$	28/17	1	0	0	4/7	-1/17
3	$x_2$	15/17	0	1	0	-3/34	5/34

Since all rows and columns are ignored, hence an optimum basic feasible solution has been reached.

$\therefore$  Optimum solution is  $x_1 = 28/17$ ,  $x_2 = 15/17$  and max.  $Z = 185/17$ .

### PROBLEM -2

$$\text{Minimum } Z = 3x_1 - 7x_2 + 5x_3$$

Subject to the constraints:

$$5x_1 - x_2 + 4x_3 \leq 15,$$

$$-3x_1 + 4x_2 \leq 8,$$

$$4x_1 + 3x_2 - 8x_3 \leq 31,$$

$$x_1, x_2, x_3 \geq 0.$$

**SOLUTION.** We have the constraints

$$5x_1 - x_2 + 4x_3 + s_1 = 15$$

$$-3x_1 + 4x_2 + s_2 = 8$$

$$4x_1 + 3x_2 - 8x_3 + s_3 = 31$$

where  $s_1, s_2, s_3$  are slack variables.

**New Simplex table**

$C_B$	ba sis	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
0	$s_1$	15	5	-1	4	1	0	0
0	$s_2$	8	-3	4	0	0	1	0
0	$s_3$	31	4	3	-8	0	0	1
-3	$x_1$	3	1	-1/5	4/5	1/5	0	0
0	$s_2$	17	0	17/5	12/5	3/5	1	0
0	$s_3$	19	0	19/5	-24/5	-4/5	0	1
-3	$x_1$	4	1	0	52/95	3/19	0	1/19
0	$s_2$	0	0	0	63/69 5	25/19	1	-17/19
7	$x_2$	5	0	1	-24/19	-4/19	0	5/19
-3	$x_1$	4	1	0	0	152/3 021	-13/15 9	755/59 47
-5	$x_3$	0	0	0	1	125/6 36	95/63 6	-85/636
7	$x_2$	5	0	1	0	-486/3 021	6/159	285/30 21

Since all rows and columns are ignored, hence an optimum basic feasible solution has been reached.

∴ Optimum solution is  $x_1 = 4, x_2 = 5, x_3 = 0$  and

max.  $Z = 23$ .

Min  $Z = -(\max Z) = -23$

**PROBLEM -3**

Maximize  $Z = 2x_1 + 3x_2$

Subject to the constraint:

$$x_1 + x_2 \leq 4,$$

$$-x_1 + x_2 \leq 1,$$

$$x_1 + 2x_2 \leq 5,$$

$$x_1, x_2 \geq 0.$$

**SOLUTION:** We have the constraints

$$x_1 + x_2 + s_1 = 4,$$

$$-x_1 + x_2 + s_2 = 1,$$

$$x_1 + 2x_2 + s_3 = 5$$

**New Simplex Table**

$C_B$	basis	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	4	1	1	1	0	0
0	$s_2$	1	-1	1	0	1	0
0	$s_3$	5	1	2	0	0	1
0	$s_1$	3/2	1/2	0	1	0	-1/2
0	$s_2$	-3/2	-3/2	0	0	1	-1/2
3	$x_2$	5/2	1/2	1	0	0	1/2
2	$x_1$	3	1	0	2	0	-1
0	$s_2$	3	0	0	3	1	-2
3	$x_2$	1	0	1	-1	0	1

Since all rows and columns are ignored, hence an opt. basic feasible solution has been reached.

∴ Optimum solution is  $x_1 = 3, x_2 = 1$  and max.  $Z = 9$ .

**PROBLEM -4**

Minimum  $Z = x_1 - 3x_2 + 2x_3$

Subject to the constraints:

$$3x_1 - x_2 + 2x_3 \leq 7,$$

$$-2x_1 + 4x_2 \leq 12,$$

$$-4x_1 + 3x_2 + 8x_3 \leq 10,$$

$$x_1, x_2, x_3 \geq 0.$$

**SOLUTION.** We have the constraints

$$3x_1 - x_2 + 2x_3 + s_1 = 7,$$

$$-2x_1 + 4x_2 + s_2 = 12,$$

$$-4x_1 + 3x_2 + 8x_3 + s_3 = 10,$$

where  $s_1, s_2, s_3$  are slack variables.

**New simplex table**

$C_B$	basiss	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
0	$s_1$	7	3	-1	2	1	0	0
0	$s_2$	12	-2	4	0	0	1	0
0	$s_3$	10	-4	3	8	0	0	1
0	$s_1$	9/2	4	-7/4	0	1	0	-1/4
0	$s_2$	12	-2	4	0	0	1	0
-2	$x_3$	5/4	-1/2	3/8	1	0	0	1/8
0	$s_1$	39/4	25/8	0	0	1	7/16	-1/4
3	$x_2$	3	-1/2	1	0	0	1/4	0
-2	$x_3$	1/8	-13/2	0	1	0	-3/32	1/8
-1	$x_1$	78/25	0	0	1	8/25	7/50	-2/25
3	$x_2$	114/25	1	0	0	4/25	8/25	-1/25
-2	$x_3$	523/25	0	1	0	52/25	643/800	79/25

Since all rows and columns are ignored, an optimum basic feasible solution has been reached.

Hence solution is

$$x_1 = 78/25, x_2 = 114/25, x_3 = 523/25 \text{ and}$$

$$\text{Min } Z = -(\text{max } Z) = -782/25$$

**III. ALTERNATIVE ALGORITHM FOR BIG-M METHOD**

To find optimal solution of any LPP by an alternative method for Big-M method, algorithm is given as follows:

Step 1. Check objective function of LPP is of maximization or minimization type. If it is to be minimization type then convert it into a maximization type by using the result:

$$\text{Min. } Z = - \text{Max. } (-Z).$$

Step 2. Check whether all  $b_i$  (RHS) are non-negative. If any  $b_i$  is negative then multiply the corresponding equation of the constraints by -1.

Step 3. Express the given LPP in standard form then obtain initial basic feasible solution.

If basic solution is non-feasible due to the constraints of the type  $\geq$  and  $=$  then we add artificial variable to the corresponding constraint in standard form. Assign very

large value  $+M$  for maximization and  $-M$  for minimization in objective function.

Step 4. Select max  $\sum x_{ij}, x_{ij} \geq 0$  for entering vector.

Step 5. Choose greatest coefficient of decision variables.

(i) If greatest coefficient is unique, then variable corresponding to this column becomes incoming variable.

(ii) If greatest coefficient is not unique, then use tie breaking technique.

Step 6. Compute the ratio with  $X_B$ . Choose minimum ratio, then variable corresponding to this row is outgoing variable. The element corresponding to incoming variable and outgoing variable becomes pivotal (leading) element.

Step 7. Use usual simplex method for this table and go to next step.

Step 8 Ignore corresponding row and column. Proceed to step 5 for remaining elements and repeat the same procedure until an optimal solution is obtain or there is an indication for unbounded solution.

Step 9: If all rows and columns are ignored, then current solution is an optimal solution.

**PROBLEM -5**

$$\text{Max } Z = 6x_1 + 4x_2$$

Subject to:

$$2x_1 + 3x_2 \leq 30,$$

$$3x_1 + x_2 \leq 24,$$

$$x_1 + x_2 \geq 3,$$

$$x_1, x_2 \geq 0.$$

**SOLUTION:** We have the constraints

$$2x_1 + 3x_2 + s_1 = 30,$$

$$3x_1 + x_2 + s_2 = 24,$$

$$x_1 + x_2 - s_3 + A_1 = 3,$$

$$x_1, x_2 \geq 0.$$

$$x_1, x_2, x_3, s_1, s_2, s_3 \geq 0.$$

where  $s_1, s_2, s_3$  are slack variables and  $A_1$  is artificial variable.

**Simplex table:**

$C_B$	basiss	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	$A_1$
0	$s_1$	30	2	3	1	0	0	0
0	$s_2$	24	3	1	0	1	0	0
-M	$A_1$	3	1	1	0	0	-1	1
0	$s_1$	14	0	7/3	1	-2/3	0	0
6	$x_1$	8	1	1/3	0	1/3	0	0
-M	$A_1$	5	0	2/3	0	-1/3	-1	0

								1
4	$x_2$	6	0	1	3/7	-2/7	0	0
6	$x_1$	6	1	0	-1/7	3/7	0	0
-M	$A_1$	-9	0	0	-2/7	-1/7	-1	1

Since all rows and column are ignored, hence an optimum solution has been reached. Therefore optimum solution is:

$$x_1 = 6, x_2 = 6 ; \text{Max} Z = 60$$

**PROBLEM- 6**

$$\text{Min } Z = 2x_1 + x_2$$

Subject to:

$$3x_1 + x_2 = 3,$$

$$4x_1 + 3x_2 \geq 6,$$

$$x_1 + x_2 \leq 3,$$

$$x_1, x_2 \geq 0.$$

**SOLUTION:** We have the constraints

$$3x_1 + x_2 + A_1 = 3,$$

$$4x_1 + 3x_2 - s_1 + A_2 = 6,$$

$$x_1 + x_2 + s_2 = 3,$$

$$x_1, x_2, s_1, s_2 \geq 0.$$

where  $s_1, s_2$  are surplus and slack variables respectively and  $A_1, A_2$  are artificial variables.

**Simplex table:**

$C_B$	basis	$x_B$	$x_1$	$x_2$	$s_1$	$A_1$	$A_2$	$s_2$
-M	$A_1$	3	3	1	0	1	0	0
-M	$A_2$	6	4	3	-1	0	1	0
0	$s_2$	3	1	1	0	0	0	1
-M	$A_1$	-3/2	0	-5/4	3/4	1	-3/4	0
-2	$x_1$	3/2	1	3/4	-1/4	0	1/4	0
0	$s_2$	3/2	0	1/4	1/4	0	-1/4	1
-1	$x_2$	6/5	0	1	-3/5	4/5	3/5	0
-2	$x_1$	3/5	1	0	-1	1	1	0
0	$s_2$	6/5	0	0	-2	3	2	1

Since all rows and column are ignored, hence an optimum solution has been reached. Therefore optimum solution is:

$$x_1 = 3/5, x_2 = 6/5 ; \text{Min } Z = 12/5$$

**IV. ALTERNATIVE ALGORITHM FOR DUAL SIMPLEX METHOD**

To find optimal solution of any LPP by an alternative method for dual simplex method, algorithm is given as follows:

Step 1. The objective function of the LPP must be maximize. If it is minimize then convert it into maximize by using the result:

$$\text{Min. } Z = - \text{Max. } (-Z).$$

Step 2. Convert all  $\geq$  constraints into  $\leq$  by multiplying the corresponding equation of the constraints by -1.

Step 3. Convert inequality constraints into equality by addition of slack variables and obtain an initial basic solution. Express the above information in the form of a table known as dual simplex table.

Step 4. Choose most negative  $X_B$ , then variable corresponding to this row becomes outgoing variable.

Select the most negative  $\sum x_{ij}$ ,  $x_{ij} \leq 0$  of, then variable corresponding to this column becomes incoming variable. The element corresponding to incoming variable and outgoing variable is pivotal (leading) element.

Step 5. Use usual simplex method for this table and go to next step.

Step 6. Ignore corresponding row and column. Proceed to step 4 for remaining elements and repeat the same procedure until an optimal solution is obtained in finite number steps or there is an indication of the non-existence of a feasible solution.

Step 7: If all rows and columns are ignored, then current solution is an optimal solution.

**PROBLEM -7**

$$\text{Minimize } Z = x_1 + 2x_2 + 3x_3$$

Subject to

$$x_1 - x_2 + x_3 \geq 4$$

$$x_1 + x_2 + 2x_3 \leq 8$$

$$x_2 - x_3 \geq 2$$

$$x_1, x_2, x_3 \geq 0$$

**SOLUTION:** We have the constraints

$$-x_1 + x_2 - x_3 + s_1 = 4$$

$$x_1 + x_2 + 2x_3 + s_2 = 8$$

$$0x_1 + x_2 + x_3 + s_3 = 2$$

$$x_1, x_2, x_3 \geq 0$$

**Initial simplex table**

$C_B$	basis	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
0	$s_1$	-4	-1	1	-1	1	0	0
0	$s_2$	8	1	1	2	0	1	0

0	$s_3$	-2	0	-1	1	0	0	1
-1	$x_1$	4	1	-1	1	-1	0	0
0	$s_2$	4	0	2	1	1	1	0
0	$s_3$	2	0	1	1	0	0	1
-1	$x_1$	6	1	0	0	-1	0	1
0	$s_2$	0	0	0	3	1	1	2
-2	$x_2$	2	0	1	-1		0	-1

Since all  $X_B$  are positive, current solution is an optimal solution.

$$x_1 = 6, x_2 = 2, x_3 = 0; \text{Max } Z^* = -10;$$

$$\text{Min } Z = 10$$

**PROBLEM -8**

Minimize  $Z = 6x_1 + 7x_2 + 3x_3 + 5x_4$

Subject to:  $5x_1 + 6x_2 - 3x_3 + 4x_4 \geq 12$

$$x_2 + 5x_3 - 6x_4 \geq 10$$

$$2x_1 + 5x_2 + x_3 + x_4 \geq 8.$$

$$x_1, x_2, x_3, x_4 \geq 0.$$

**SOLUTION:** We have the constraints

$$-5x_1 - 6x_2 + 3x_3 - 4x_4 + s_1 = -12$$

$$-x_2 - 5x_3 + 6x_4 + s_2 = -10$$

$$-2x_1 - 5x_2 - x_3 - x_4 + s_3 = -8.$$

$$x_1, x_2, x_3, s_1, s_2, s_3 \geq 0.$$

where  $s_1, s_2, s_3$  are slack variables.

**Simplex table:**

Since all  $X_B$  are positive, current solution is an optimal solution.

$$x_1 = 0, x_2 = \frac{30}{11}, x_3 = \frac{16}{11}, x_4 = 0. \text{Min. } Z = \frac{258}{11}.$$

**PROBLEM -9**

Maximize  $Z = x_1 + 2x_2$

Subject to:  $3x_1 + x_2 \geq 3$

$$4x_1 + 3x_2 \geq 6$$

$$x_1 + 2x_2 \geq 3,$$

$$x_1, x_2 \geq 0$$

**SOLUTION:** We have the constraints

$$-3x_1 - x_2 + s_1 = -3$$

$$-4x_1 - 3x_2 + s_2 = -6$$

$$-x_1 - 2x_2 + s_3 = -3$$

$$x_1, x_2, x_3 \geq 0.$$

$$x_1, x_2, x_3, s_1, s_2, s_3 \geq 0.$$

**Simplex table:**

$C_B$	basis	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$
0	$s_1$	-3	-3	-1	1	0	0
0	$s_2$	-6	-4	-3	0	1	0
0	$s_3$	-3	-1	-2	0	0	1
0	$s_1$	3/2	0	5/4	1	-3/4	0
1	$x_1$	3/2	1	3/4	0	-1/4	0
0	$s_3$	5/4	0	5/4	0	-1/4	1
0	$s_1$	0	0	0	1	-1	1
1	$x_1$	3/5	1	0	0	-2/5	3/5
2	$x_2$	6/5	0	1	0	1/5	-4/5

Since all  $X_B$  are positive, current solution is an optimal solution.

$$x_1 = 3/5, x_2 = 6/5 \text{ MAX } Z = 3$$

**V. CONCLUSION**

An alternative methods for simplex method, Big M method and dual simplex method have been derived to obtain the solution of linear programming problem. The proposed algorithms have simplicity and ease of understanding. This reduces number of iterations and improves the optimum solutions in most of the cases. These methods save valuable time as there is no need to calculate the net evaluation  $Z_j - C_j$ .

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